
Distributed Video Systems

Chapter 3

Storage Technologies

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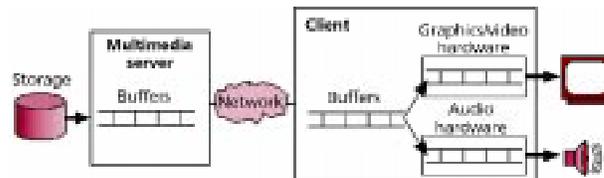
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3.1 Introduction

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- System Model



- Challenges

- ◆ Real-time storage and retrieval:
 - Continuous media data must be presented using the same timing sequence with which they were captured.
 - Any deviation from this timing sequence can lead to artifacts such as jerkiness in video motion, pops in audio, or possibly complete unintelligibility.

3.1 Introduction

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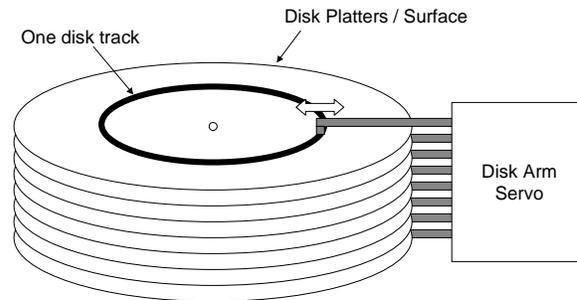
- Challenges

- ◆ Real-time storage and retrieval:
 - Media components may also need *synchronization*. For example, a video stream must synchronize an audio stream in a movie.
- ◆ High data transfer rate and large storage space:
 - Digital video and audio playback demands a high data transfer rate, so that storage space is rapidly filled. (E.g. MPEG1 ~ 1.5Mbps, MPEG2 ~ 4Mbps)
 - The server must efficiently store, retrieve, and manipulate data in large quantities at high speeds.

3.2 Magnetic Disks

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- Disk Model



- ◆ The disk platters spin at speed from 3600rpm to 10000rpm;
- ◆ Disk heads in all platters move together.
- ◆ A disk track is further divided into disk sectors.

3.2 Magnetic Disks

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- Disk Model

- ◆ Fixed Delays
 - Processing delay at disk controller;
 - Delay at data bus (e.g. SCSI) between disk and controller;
 - Head-switching time;
- ◆ Variable Delays
 - Rotational Latency
 - Depends on position and spindle speed
 - Seek time
 - Depends on number of tracks to seek
 - Transfer Time
 - Depends on how much data to transfer to host

3.2 Magnetic Disks

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- Disk Model

- ♦ Disk-Seek Time Function:

$$T_{seek}(n) = \alpha + \beta\sqrt{n}$$

Number of tracks to seek
Seek-time constant (sec)
Fixed overhead (sec)

- ♦ Total Disk-Read Time Function:

$$T_{read}(n) = \alpha + \beta\sqrt{n} + T_{latency} + \frac{Q}{R_{disk}}$$

Size of data to read (Bytes)
Disk transfer rate (Bytes/sec)
Rotational latency (sec)

3.2 Magnetic Disks

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- Typical Disk Parameters

- ♦ Seagate 4GB ST12400N (SCSI-2)

Disk Parameter	Value
Spindle speed	5411 rpm
Max latency (r)	11ms
Number of tracks	2621
Raw transfer rate	3.35MB/s
Single-track seek	1ms
Max full-stroke seek	19ms

3.2 Magnetic Disks

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- Typical Disk Parameters
 - ♦ SCSI Variants

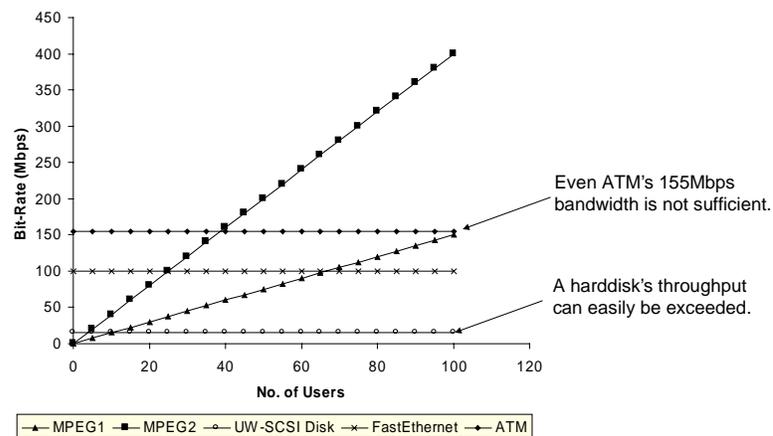
Types	Variants	Max. Speed	Number of Devices	Max. Cable Length
SCSI-1	-	5 MB/s	8	6m
SCSI-2	Fast SCSI	10 MB/s	8	1.5m~3m
	Fast Wide SCSI	20 MB/s	16	1.5m~3m
SCSI-3	Ultra SCSI	20 MB/s	8	1.5m
	Wide Ultra SCSI	40 MB/s	16	1.5m
	Ultra2 SCSI	40 MB/s	8	12m
	Wide Ultra2 SCSI	80 MB/s	16	12m
	Ultra3 SCSI	80 MB/s	8	12m
	Wide Ultra3 SCSI	160 MB/s	16	12m
Fibre Channel	FC-AL	100~200MB/s	126	30m~10km

- Note that the "Max. Speed" is the top speed of the interface.
- The actual achievable speed depends on the performance of the connected disks.

3.3 Video Retrieval

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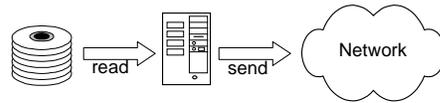
- The Bandwidth Landscape:



3.3 Video Retrieval

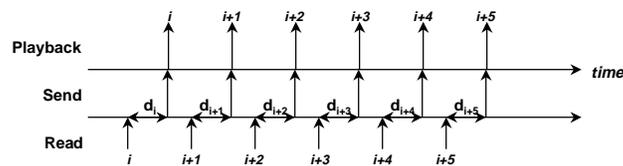
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- Single-Stream Retrieval



- ◆ Ideal Disk (Constant Service Time)

Constant delay: $d_i = d_j \forall i, j$



Assumes zero transmission time in network.

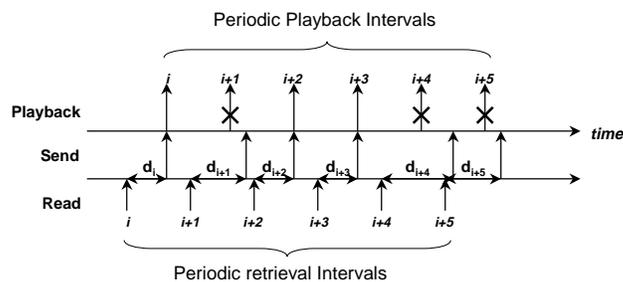
3.3 Video Retrieval

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- Single-Stream Retrieval

- ◆ In Practice (Variable Service Time)

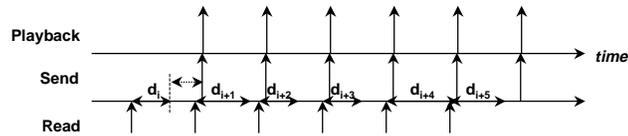
- Variable delay can cause playback glitches:



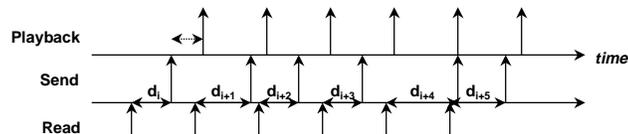
3.3 Video Retrieval

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- Single-Stream Retrieval
 - ♦ In Practice (Variable Service Time)
 - Buffering At Server:



- Buffering At Receiver:



3.3 Video Retrieval

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- Multi-Stream Retrieval
 - ♦ One Disk Per Stream
 - Simple but wasteful because disk bandwidth is usually much larger than video bit-rate.
 - E.g. >10Mbps for HD, but MPEG2 only ~4Mbps.
 - ♦ Multiple Streams Per Disk
 - A *disk scheduling algorithm* is required to ensure that the individual streams will not interfere with each other, and the delay constraint is met.
 - There are many disk scheduling algorithms, each with its own strengths and weaknesses.

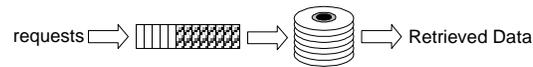
3.4 Disk Scheduling

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- Conventional Disk Scheduling Algorithms

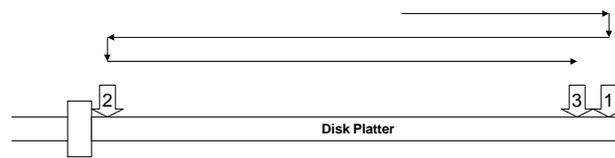
- First-Come-First-Serve (FCFS)

- Service requests in the order they arrive.



- Simple but poor disk utilization.

– Example:



Very long seek time in this example.

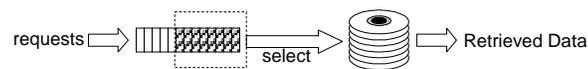
3.4 Disk Scheduling

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- Conventional Disk Scheduling Algorithms

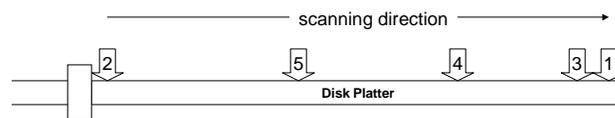
- SCAN

- Service requests along *scanning* direction.



- Better disk utilization but potentially long *round time*.

– Example:



Service Order: 2 5 4 3 1

Note request 1 has to wait longer even it arrives first!

3.4 Disk Scheduling

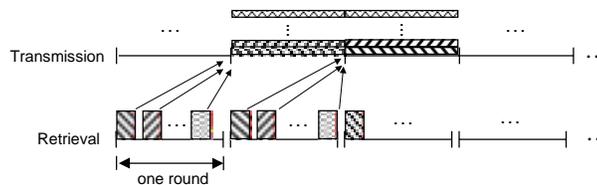
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- Multimedia Disk Scheduling Algorithms
 - ♦ Earliest Deadline First (EDF)
 - This algorithm schedules the media block with the earliest deadline for retrieval.
 - Likely to yield *excessive* seek time and rotational latency, and *poor* server-resource utilization can be expected.
 - ♦ Scan-EDF
 - Same as EDF except using SCAN to schedule requests having the same deadline.

3.4 Disk Scheduling

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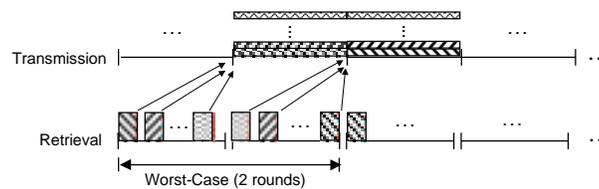
- Disk Scheduling Algorithms for VoD Servers
 - ♦ Characteristic of Continuous Media
 - Periodic retrieval of fixed-size data blocks;
 - The entire retrieval schedule is known beforehand.
 - ♦ Round-Based Disk Scheduling
 - Read one data block for each video stream in each round.
 - Retrievals in a round are serviced using SCAN.



3.4 Disk Scheduling

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- Disk Scheduling Algorithms for VoD Servers
 - ♦ Round-Based Disk Scheduling
 - To ensure the continuity of data flow for transmission, we need **two buffers per video stream**.
 - Limitations
 - All video streams must have the same data rate; or
 - The data rate must be an integer multiple of a base data rate.



3.5 Admission Control

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- Admission Control
 - ♦ Motivation
 - A VoD system only have finite capacity. Hence a mechanism must be used to admit and reject users to avoid system overload.
 - ♦ Types of Admission Control Algorithms
 - Deterministic
 - Worst-case scenarios are used to guarantee the service of existing users.
 - Statistical
 - Statistical behaviors of the system are used to provide *probabilistic* guarantee. E.g. meeting deadline 99% of the time.
 - Observational
 - Current system status like utilizations are used to evaluate the admission of new users.

3.5 Admission Control

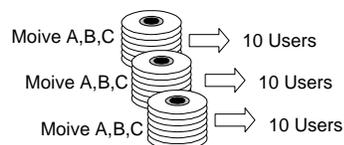
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- Dealing with Missed Deadlines
 - ◆ Why?
 - Deadlines could be missed if the admission control algorithm is statistical or some other unexpected events occur.
 - ◆ What to do?
 - Ignore It
 - Causes service degradations such as jerky video, decoding error, scrambled video, audio clicks, etc.
 - Depends on how much and what kind of data is missed.
 - Error Concealment
 - Repeating data (previous frame, audio packet, etc.)
 - Skipping video frame
 - Lower the resolution (temporary)

3.6 I/O Bandwidth

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- Increasing Disk Throughput
 - ◆ Background
 - A single disk's through can serve a very limited number of concurrent users.
 - For example, a SCSI harddisk can serve around 10 MPEG1 video streams and 3~4 MPEG2 video streams.
 - ◆ Replication
 - Use multiple disks, each carry a separate copy of a movie.
 - Expensive since movie is large in size.



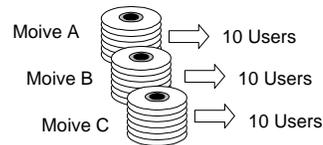
3.6 I/O Bandwidth

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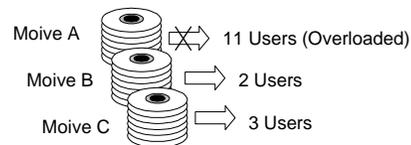
- Increasing Disk Throughput

- ◆ Partition

- Use multiple disks, each carry different movie titles.



- Same total storage but poor load-balancing.



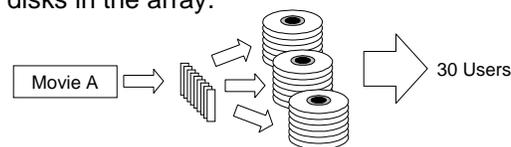
3.6 I/O Bandwidth

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- Increasing Disk Throughput

- ◆ Disk Striping (Disk Array)

- Divides a video stream into units and distributes over all disks in the array.



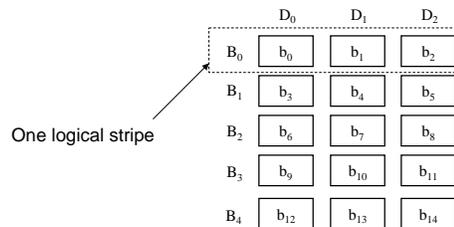
Data Organization:

	D ₀	D ₁	D ₂
B ₀	b ₀	b ₁	b ₂
B ₁	b ₃	b ₄	b ₅
B ₂	b ₆	b ₇	b ₈
B ₃	b ₉	b ₁₀	b ₁₁
B ₄	b ₁₂	b ₁₃	b ₁₄

3.6 I/O Bandwidth

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- Increasing Disk Throughput
 - ♦ Disk Striping (Disk Array)
 - One logical *stripe* is retrieved per stream per round.

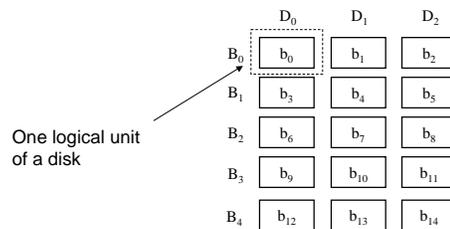


- Hence the throughput is N times those of a single disk if there are N disks in the array.
- The disks are *spindle synchronized*.

3.6 I/O Bandwidth

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- Increasing Disk Throughput
 - ♦ Disk Interleaving
 - Same as disk striping except one logical unit is retrieved from one of the disk per stream per round.



- Hence each disk can serve a different stream at the same time, or multiple streams are served concurrently.
- The disks are not spindle synchronized and operates independently.

3.7 Storage Capacity

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- Tertiary Storage and Storage Hierarchies

- ♦ Motivation

- While magnetic disks are suitable for use in VoD systems due to the high throughput and low latency, *they are still expensive.*
- For applications like video library where large number of videos must be archived, storing all video in disks will become prohibitively expensive (and unnecessary).

- ♦ Tertiary Storage

Feature	Magnetic Disk	Optical Disk	Low-end Tape	High-end Tape
Capacity	9GB	200GB	500GB	10TB
Mount time	None	20 secs	60 secs	90 secs
Transfer Rate	2MBps	300KBps	100KBps	1MBps
Cost	\$5,000	\$50,000	\$50,000	\$0.5M to \$1M
Cost/GB	\$555	\$125	\$100	\$50

3.7 Storage Capacity

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- Tertiary Storage and Storage Hierarchies

- ♦ Tertiary Storage

- Pros
 - Removable media like optical disks and tapes are less expensive in terms of cost per GB.
- Cons
 - Lower data transfer rate;
 - Very long random access time.

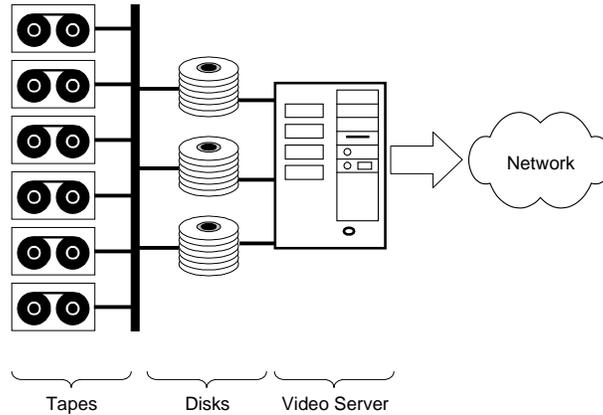
- ♦ Storage Hierarchy

- Combines the cost-effectiveness of tertiary storage with the performance of magnetic disks.
- Tertiary storage are used for permanent storage and the magnetic disks used as a cache for video delivery.

3.7 Storage Capacity

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- Tertiary Storage and Storage Hierarchies
 - ♦ Storage Hierarchy



3.7 Storage Capacity

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- Tertiary Storage and Storage Hierarchies

- ♦ Storage Hierarchy

- Scheme 1:

- Store the beginning segments of videos in magnetic disk and the rest in tertiary storage;
- Starts delivery from magnetic disk while downloading the rest of the video from the tertiary storage.

- Scheme 2:

- Downloads an entire video from tertiary storage to magnetic disks for delivery.
- Manage the disk storage using most-recently-used policy.
- Long startup time for uncached video but the caching should perform well since only a small number of video will be popular at any one time.